

ON THE PROBABILITY OF CORRECT SELECTION
IN THE SUBSET SELECTION PROBLEM

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I. Introduction.

Subset selection deals with the problem of selecting a random non-empty subset of populations out of say, k populations, with the aim that the selected populations are "close" in some sense to the best population. In particular, a subset including the best population is called a correct selection (CS). The classical condition on subset selection procedures is to require $P(\text{CS}) \geq P^*$. Usually P^* is chosen to be greater than $1/k$.

There seems to be some confusion as to why and whether $1/k$ is an appropriate lower bound for P^* . Gupta (1965) states that one should choose $P^* \geq 1/k$ because for $P^* < 1/k$ there always exist a no-data decision rule. Gibbons, Olkin and Sobel (1977) states that $1/2^k$ is the appropriate lower bound, but the justification given for this bound is incorrect. The bound $1/2^k$ is claimed to be obtained by the rule that selects each population with probability $1/2$. However, this is not a subset selection rule since it may select an empty set. Also, $P(\text{CS}) = 1/2$ for this procedure. Bechhofer and Santner (1979) support the lower bound $1/k$ on the basis of certain minimax arguments for no-data decision-rules.

The aim of this note is to clarify this issue. This author thinks that a lower bound on P^* should depend only on the decision-space and

the class of procedures under consideration. From this point of view it is shown in Section 1 that for some reasonable classes, $1/k$ is the correct lower bound, in the sense that no procedure in those classes can achieve a P^* less than $1/k$. However, it turns out that for several classes $1/k$ is not the appropriate bound. It is also shown that a procedure with $P(CS) < 1/k$ has certain undesirable properties, which gives an argument in favor of considering only classes of procedures that has $1/k$ as the lower bound.

Section 2 deals with no-data rules. It is shown that Gupta's statement is incorrect if one only considers permutation-invariant no-data procedures. Bechhofer and Santner's approach is also briefly discussed.

2. $P(CS)$ for Monotone, Ordered and Permutation-Invariant Procedures.

We shall consider the following situation. The k populations are denoted by π_1, \dots, π_k . π_i is characterized by a real-valued parameter θ_i . X_i is the observation from π_i . X_1, \dots, X_k are assumed to be independent. $F_{\theta_i}(x)$ is the distribution function of X_i . It is assumed that $F_{\theta}(x)$ is stochastically increasing in θ , and continuous in x for each θ in the parameter-space $\Theta \subset \mathbb{R}$. This class of distribution functions is denoted by \mathcal{F}_C . Let $\Omega = \Theta^k$. Let $\theta_{[1]} \leq \dots \leq \theta_{[k]}, X_{[1]} \leq \dots \leq X_{[k]}$ denote the ordered θ_i 's and X_i 's. $X_{(i)}, \pi_{(i)}$ correspond to $\theta_{[i]}$. $\pi_{(k)}$ is defined to be the best population. The decision-space \mathcal{A} is the set of all non-empty subsets of (π_1, \dots, π_k) . A subset selection-rule δ is for each observed $\underline{x} = (x_1, \dots, x_k)$ a probability-measure $\delta(a|\underline{x})$ over $a \in \mathcal{A}$. For a procedure δ , the individual selection probabilities are given by:

$$\psi_i^\delta(\underline{x}) = P(\text{selecting } \pi_i \text{ with rule } \delta|\underline{x}) = \sum_{a \ni i} \delta(a|\underline{x})$$

Let $\psi(i)^\delta$ correspond to $0_{[i]}$. Then the classical P^* condition is

$$\inf_{0 \in \Omega} P_0(CS|\delta) = \inf_{0 \in \Omega} E_{\theta \psi(k)^\delta}(\underline{X}) = P^* \quad (1)$$

Let us for convenience denote $\inf_{\theta \in \Omega} P_0(CS|\delta)$ by $P^*(\delta)$. The range of possible values of P^* will depend upon the class \mathcal{L} under consideration.

Suppose $\inf_{\delta \in \mathcal{L}} P^*(\delta) = \alpha$, then the principle is that P^* should be at least α , because no rules in \mathcal{L} can achieve a P^* less than α . Consider for example the class \mathcal{L}_G of Gupta's rules (see Gupta (1965)).

$$\psi_i = 1 \text{ iff } X_i > X_{[k]}^{-d}, \quad d > 0.$$

Here $\inf_{\delta \in \mathcal{L}_G} P^*(\delta)$ is $1/k$, achieved by the rule corresponding to $d=0$.

Hence for \mathcal{L}_G , P^* should be at least $1/k$.

This principle will be applied to different classes of procedures to find out if $1/k$ is the natural lower bound. In order to define the class \mathcal{L}_I of permutation-invariant procedures, let g be a permutation of $(1, \dots, k)$ such that g_i is the new position of element i under permutation g . Then $g\underline{x}$ is defined by $(g\underline{x})_i = x_{g^{-1}i}$, and $g_a = \{g_i : i \in a\}$ for $a \in \mathcal{A}$.

Definition 1. $\delta \in \mathcal{L}_I$ if for each permutation g

$$\delta(g_a | g\underline{x}) = \delta(a | \underline{x}) \quad \forall a \in \mathcal{A}, \forall \underline{x}.$$

Definition 2. δ is said to be ordered if

$$\psi_i^\delta(\underline{x}) \leq \psi_j^\delta(\underline{x}) \quad \text{when } x_i < x_j$$

\mathcal{L}_0 denotes the class of ordered procedures.

Definition 3. δ is called monotone if for each i ,

$$\psi_i^\delta(\underline{x}) \leq \psi_i^\delta(\underline{y}) \quad \text{if} \quad x_i \leq y_i \quad \text{and} \quad x_j \geq y_j \quad \forall j \neq i.$$

Let \mathcal{S}_M be the class of monotone procedures.

\mathcal{S}_I , \mathcal{S}_O , \mathcal{S}_M are the three basic classes of procedures we consider.

If \mathcal{S} is one of the three or a combination of these, the basic question to answer is whether or not the following statement is true:

$$\inf_{\delta \in \mathcal{S}} P^*(\delta) = 1/k \quad (2)$$

Let us first discuss the relationship between the three classes. Clearly a rule can be monotone and not ordered or vice versa. The following results also hold.

Lemma 1.

- (i) $\delta \in \mathcal{S}_I \not\Rightarrow \delta \in \mathcal{S}_O$ (i.e. $\mathcal{S}_I - \mathcal{S}_O$ is non-empty.)
- (ii) $\delta \in \mathcal{S}_O \not\Rightarrow \delta \in \mathcal{S}_I$ (i.e. $\mathcal{S}_O - \mathcal{S}_I$ is non-empty.)
- (iii) $\mathcal{S}_{I,M} \subset \mathcal{S}_{I,O}$ and $\mathcal{S}_{I,O} - \mathcal{S}_{I,M}$ is non-empty.

Here $\mathcal{S}_{I,M} = \mathcal{S}_I \cap \mathcal{S}_M$ and $\mathcal{S}_{I,O} = \mathcal{S}_I \cap \mathcal{S}_O$.

Proof.

(i) is obvious. E.g. the rule that selects π_i if and only if $X_i = X_{[k-1]}$ is permutation-invariant but not ordered.

(ii). Consider the following rule δ :

If $X_1 = X_{[k]}$: select π_1 .

If $X_1 < X_{[k]}$: select $\pi_{[k]}$, $\pi_{[k-1]}$, where $\pi_{[i]}$ corresponds to $X_{[i]}$.

$\delta \in \mathcal{S}_O$, but $\delta \notin \mathcal{S}_I$.

(iii). First we note that $\delta \in \mathcal{L}_I$ implies that $\psi_i^\delta(\underline{x}) = \psi_{gi}^\delta(g\underline{x}) \quad \forall (g, i, \underline{x})$.

Assume $\delta \in \mathcal{L}_{I, M}$. Let \underline{x} be such that $x_i < x_j$.

We shall show that $\psi_i^\delta(\underline{x}) \leq \psi_j^\delta(\underline{x})$.

Let g be the permutation with $gi=j, gj=i, gl=l \quad \forall l \neq i, j$, and let $y=g\underline{x}$.

Then $y_i = x_j > x_i$ and $y_l \leq x_l \quad \forall l \neq i$. Hence $\psi_i^\delta(\underline{x}) \leq \psi_i^\delta(\underline{y})$ from Definition 3,

and $\psi_j^\delta(\underline{x}) = \psi_{gj}^\delta(g\underline{x}) = \psi_i^\delta(\underline{y}) \geq \psi_i^\delta(\underline{x})$, which proves the first statement. Let now

$k > 3$, and consider the following rule.

$$\delta: \text{select } \pi_i \Leftrightarrow X_i \geq \min(X_{[k]}, \frac{X_{[k]}}{\bar{X}}) ; \bar{X} = \frac{1}{k} \sum_{i=1}^k X_i .$$

δ is clearly in $\mathcal{L}_{I, 0}$. We shall show that δ is not monotone.

Let $x_i = 3/2$ for $i \leq k-1$ and $x_k = 2$, and let $y_i = 0$ for $i \leq k-2$, $y_{k-1} = 3/2$, $y_k = 2$.

Here $x_{k-1} = y_{k-1}$ and $y_j \leq x_j \quad \forall j \neq k-1$. It is readily seen that $\psi_{k-1}(\underline{x}) = 1$

and $\psi_{k-1}(\underline{y}) = 0$.

Q.E.D.

The results about $P^*(\delta)$ for the classes $\mathcal{L}_I, \mathcal{L}_0, \mathcal{L}_M$ are given in the following

Theorem 1.

(a) $\inf_{\delta \in \mathcal{L}_I} P^*(\delta) = 0$, provided $F_\theta(x) \rightarrow 1$ as $\theta \rightarrow \inf \Theta$.

(b) $\inf_{\delta \in \mathcal{L}_0} P^*(\delta) = (1/k)^2$; $\forall F_\theta \in \mathcal{F}_C$.

(c) $\inf_{\delta \in \mathcal{L}_M} P^*(\delta) = 0$; $\forall F_\theta \in \mathcal{F}_C$.

Proof.

(a) Consider the rule

$$\delta: \text{select } \pi_i \Leftrightarrow X_i = X_{[1]} .$$

$\delta \in \mathcal{L}_I$ so we may assume $\theta_i = \theta_{[i]}$.

Then

$$\begin{aligned} P_{\underline{\theta}}(CS|\delta) &= P(X_k \leq X_i ; \forall i \leq k-1) = \int \prod_{i=1}^{k-1} P(X_i \geq x) dF_{\theta_k}(x) \\ &= \int \prod_{j=1}^{k-1} (1 - F_{\theta_j}(x)) dF_{\theta_k}(x). \end{aligned}$$

Let $\theta_j \rightarrow \inf \Theta \quad \forall j \leq k-1$ and keep θ_k fixed. Then $F_{\theta_j}(x) \rightarrow 1, \forall x$ and from Lebesgue's convergence theorem $P_{\underline{\theta}}(CS|\delta) \rightarrow 0$. Hence $P^*(\delta) = 0$.

(b) We observe that $a \neq \phi$ implies $\sum_{a \in \mathcal{A}} \delta(a|\underline{x}) = 1 \quad \forall \underline{x}$ and hence

$$\sum_{i=1}^k \psi_i^\delta(\underline{x}) \geq \sum_{a \in \mathcal{A}} \delta(a|\underline{x}) = 1. \quad (3)$$

Let $\delta \in \mathcal{S}_0$. Then $x_i = x_{[k]} \Rightarrow \psi_i(\underline{x}) \geq 1/k$ from (3)

Hence:

$$\begin{aligned} P_{\underline{\theta}}(CS|\delta) &= E_{\underline{\theta}}(\psi(k)|\delta) \geq \int \psi(k)(\underline{x}) \prod_{i=1}^k dF_{\theta_i}(x_i) \geq \frac{1}{k} P_{\underline{\theta}}(X(k) = X_{[k]}) \\ &\quad \{ \underline{x} : x_{(k)} = x_{[k]} \} \end{aligned}$$

Now $\inf_{\underline{\theta} \in \Omega} P_{\underline{\theta}}(X(k) = X_{[k]}) = P_{\theta_1 = \dots = \theta_k} (X(k) = X_{[k]}) = 1/k$,

since F_{θ} is stochastically increasing in θ .

Hence: $P_{\underline{\theta}}(CS|\delta) \geq (1/k)^2 \quad \forall \underline{\theta} \in \Omega$.

The lower bound is achieved by the following rule in \mathcal{S}_0 :

If $x_i = x_{[k]}$ and $i < k$: $\delta(a_i|\underline{x}) = 1$

If $x_k = x_{[k]}$: $\delta(a_i|\underline{x}) = 1/k$ for $i=1, \dots, k$.

Here $a_i = \{\pi_i\}$.

(c) is obvious e.g. let δ be given by:

$$\psi_i^\delta(\underline{x}) = 1 \Leftrightarrow x_i \geq x_{[k-1]} \quad \text{for } i: k-1$$

$$\psi_k^\delta(\underline{x}) = 0 \quad \forall \underline{x}$$

$\delta \in \mathcal{D}_M$ and if $0_k = \max_{1 \leq i \leq k-1} 0_i$ then $P_{\underline{0}}(CS|\delta) = 0$. Q.E.D.

From Theorem 1, we see that none of the three properties, permutation-invariant, monotone or ordered, alone insures (2).

A desirable property of a procedure δ is unbiasedness.

Definition 4. δ is said to be unbiased if

$$i < j \Rightarrow E_{\underline{\theta}} \psi^{\delta}(i) \leq E_{\underline{\theta}} \psi^{\delta}(j)$$

(Some authors, e.g. Gupta (1965) and Nagel (1970) use the terminology "monotone" for this property.)

Let S be the size of the selected subset. Then

$$E_{\underline{\theta}}(S|\delta) = \sum_{i=1}^k E_{\underline{\theta}} \psi^{\delta}(i)$$

Lemma 2.

$P^*(\delta) < 1/k \Rightarrow \delta$ is not unbiased.

Proof. Let $\underline{0} \in \Omega$, arbitrary, and assume that δ is unbiased. Then

$$k \cdot P_{\underline{0}}(CS|\delta) = k E_{\underline{0}} \psi^{\delta}(k) \geq \sum_{i=1}^k E_{\underline{0}} \psi^{\delta}(i) = E_{\underline{0}}(S|\delta) \geq 1, \text{ since } S \geq 1.$$

Hence $P^*(\delta) \geq 1/k$.

Q.E.D.

So, if one only wants to consider unbiased procedures at least (2) must be satisfied. From Theorem 1 we see that there are biased procedures in each of the classes \mathcal{D}_I , \mathcal{D}_O , \mathcal{D}_M . It turns out (see Theorem 2 below) that stronger results can be obtained for some combinations of the three classes. Also restricting attention to non-randomized procedures can lead to different results. δ is non-randomized if for each x there exists $a \in \mathcal{A}$, such that $\delta(a|\underline{x}) = 1$. For a given class \mathcal{D} , let \mathcal{D}^n denote the class of non-randomized procedures in \mathcal{D} . E.g. \mathcal{D}_O^n is the

class of non-randomized ordered procedures. Our basic question is now answered by the following results.

Theorem 2.

- (a) $\inf_{\delta \in \mathcal{D}_{I,M}} P^*(\delta) = 1/k, \quad \forall F_\theta \in \mathcal{F}_C.$
- (b) $\inf_{\delta \in \mathcal{D}_0^n} P^*(\delta) = \inf_{\delta \in \mathcal{D}_{I,0}^n} P^*(\delta) = \inf_{\delta \in \mathcal{D}_{M,0}^n} P^*(\delta) = 1/k; \quad \forall F_\theta \in \mathcal{F}_C.$

Here $\mathcal{D}_{M,0} = \mathcal{D}_M \cap \mathcal{D}_0.$

- (c) $\inf_{\delta \in \mathcal{D}_{I,0}} P^*(\delta) < 1/k$ for some $F_\theta \in \mathcal{F}_C$
- (d) $\inf_{\delta \in \mathcal{D}_{M,0}} P^*(\delta) \leq 1/2k, \quad \forall F_\theta \in \mathcal{F}_C.$

Proof.

(a) Let $\delta \in \mathcal{D}_{I,M}$. Then from Nagel (1970):

$$P^*(\delta) = \inf_{\underline{\theta} \in \Omega_0} P_{\underline{\theta}}(CS|\delta), \text{ where } \Omega_0 = \{\underline{\theta} : \theta_1 = \dots = \theta_k\}.$$

For $\underline{\theta} \in \Omega_0$: $E_{\underline{\theta}}(\psi_1^\delta) = \dots = E_{\underline{\theta}}(\psi_k^\delta) = P_{\underline{\theta}}(CS|\delta)$, since $\delta \in \mathcal{D}_I$.

Hence: $E_{\underline{\theta}}(S|\delta) = kP_{\underline{\theta}}(CS|\delta)$. Since $S \geq 1$, it follows that $\inf_{\mathcal{D}_{I,M}} P^*(\delta) \geq 1/k$.

The lower bound is obtained by the rule that selects π_i if and only if $X_i = X_{[k]}$.

(b) Let now $\delta \in \mathcal{D}_0^n$. Then $X_{(k)} = X_{[k]} \Rightarrow X_{(k)} > X_{(i)} \quad \forall i \leq k-1$ with F_θ -probability 1, which implies $\psi_{(k)}(X) = 1$ with probability 1. Hence

$$P_{\underline{\theta}}(CS|\delta) = E_{\underline{\theta}}\psi_{(k)}^\delta \geq P_{\underline{\theta}}(X_{(k)} = X_{[k]}) \geq 1/k.$$

The lower bound is attained by the same rule as in (a).

(c) Let $F_{\theta} \in \mathcal{F}_C$ be such that there exists $0' < 0''$, $a < b$ for which

$$F_{\theta'}(a) = .49 \quad 1 - F_{\theta'}(b) = .5$$

$$F_{\theta''}(a) = .01 \quad 1 - F_{\theta''}(b) = .5$$

Let $\theta_1 = \dots = \theta_{k-1} = \theta'$, $\theta_k = \theta''$.

Consider the following rule.

$$\delta(a_i | \underline{x}) = 1 \quad \text{iff} \quad x_i > b \text{ and } a < x_j \leq b \quad \forall j \neq i$$

for $i=1, \dots, k$

$$\delta(a_i | \underline{x}) = \dots = \delta(a_k | \underline{x}) = 1/k, \text{ otherwise.}$$

It is readily seen that $\delta \in \mathcal{A}_{I,0}$.

$$P_{\theta}(\text{CS} | \delta) = E_{\theta} \psi_k^{\delta} = P(a < X_i \leq b, \forall i \leq k-1 \text{ and } X_k > b) \\ + \frac{1}{k} (1 - P(A))$$

Here:

$$P(A) = P\left[\bigcup_{i=1}^k (X_i > b \text{ and } a < X_j \leq b \quad \forall j \neq i)\right] = (.01)^{k-2} \frac{k-1}{2} + (.01)^{k-1} \cdot \frac{1}{2}$$

This gives:

$$E_{\theta} \psi_k^{\delta} = \frac{1}{k} - \left(\frac{P(A)}{k} - \frac{1}{2} (.01)^{k-1}\right) \\ = \frac{1}{k} - (.24) (.01)^{k-2} \left(\frac{k-1}{k}\right) < \frac{1}{k}$$

(d) Consider the following procedure δ given by:

$$\text{If } x_i = x_{[k]} \text{ and } i < k, \text{ then } \delta(a_i | \underline{x}) = 1$$

$$\text{If } x_k = x_{[k]}, \text{ then } \delta(a_{[k-1]} | \underline{x}) = \delta(a_{[k]} | \underline{x}) = \frac{1}{2}.$$

Here $a_{[i]} = \{a_{[i]}\}$.

$\delta \in \mathcal{D}_{M,0}$. Let $\Omega_k = \{\underline{\theta} : \theta_k = \theta_{[k]}\}$.

Then

$$\inf_{\underline{\theta} \in \Omega_k} P_{\underline{\theta}}(\text{CS}|\delta) = \inf_{\underline{\theta} \in \Omega_k} E_{\underline{\theta}} \psi_k^\delta = \frac{1}{2} \inf_{\underline{\theta} \in \Omega_k} P(X_k = X_{[k]}) = \frac{1}{2k}.$$

Q.E.D.

Remarks.

1. From the proofs of Theorem 1 (a), (c) and Theorem 2(a) we see that for the classes \mathcal{D}_I , \mathcal{D}_M and $\mathcal{D}_{I,M}$ the same results hold when restricting attention to non-randomized procedures.
2. It does not necessarily follow from (2) that all $\delta \in \mathcal{D}$ are unbiased. However, for the class $\mathcal{D}_{I,M}$, Nagel (1970) showed that all $\delta \in \mathcal{D}_{I,M}$ are also unbiased.
3. Since $\mathcal{D}_{I,M} \subset \mathcal{D}_{I,0}$ we see from Theorem 2(a), that it is essentially required that a procedure is permutation-invariant, ordered and monotone for (2) to hold, although for non-randomized procedures it is enough that the procedure is ordered.

We conclude this section with a few observations about the discrete distribution-case. Let \mathcal{F}_d be the class of all stochastically increasing discrete $F_\theta(x)$, $\theta \in \Theta \subset \mathbb{R}$. The results for \mathcal{D}_I , $\mathcal{D}_{I,M}$, $\mathcal{D}_{I,0}$, \mathcal{D}_M are essentially the same as before. It can also be shown that for the classes \mathcal{D}_0^n , \mathcal{D}_0 , $\mathcal{D}_{M,0}^n$, $\mathcal{D}_{M,0}$ we now get that $\inf P^*(\delta) = 0$ for some $F_\theta \in \mathcal{F}_d$. This differs from the results for \mathcal{D}_C .

3. $P(\text{CS})$ for Permutation-Invariant "No-Data" Rules.

By definition, δ is a no-data rule if it is independent of x , i.e.

$$\delta(a|x) = \delta(a) \quad \forall a \in \mathcal{A}, \forall x, \text{ so that}$$

$$\sum_{a \in \mathcal{A}} \delta(a) = 1.$$

δ is permutation-invariant if $\delta(ga) = \delta(a)$, $\forall g, \forall a$.

If $|a|$ denotes the size of a , then

$$\delta(ga) = \delta(a) \quad \forall g, \forall a \in \mathcal{A} \Leftrightarrow \delta(a) = \delta(a') \quad \text{if } |a| = |a'|$$

Let p_i be the probability that a subset of size i selected, i.e.

$$p_i = \sum_{\{a: |a|=i\}} \delta(a) = \binom{k}{i} \delta(\{1, \dots, i\}), \quad \text{if } \delta \text{ is permutation-invariant.}$$

Let $\underline{p} = (p_1, \dots, p_k)$. \underline{p} characterizes a permutation-invariant no-data rule, since for any a with size i

$$\delta(a) = p_i / \binom{k}{i}.$$

One way to select according to this rule in practice is first to select a subset size according to \underline{p} . Then given size i , one chooses a randomly, i.e. each subset of size i have probability $\binom{k}{i}^{-1}$ of being selected. It is readily seen (also shown by Bechhofer and Santner (1979)) that

$$P_{\underline{\theta}}(CS|\underline{p}) = \sum_{i=1}^k p_i \frac{i}{k}; \quad \text{independent of } \underline{\theta}$$

and

$$E_{\underline{\theta}}(S|\underline{p}) = \sum_{i=1}^k i p_i; \quad \text{independent of } \underline{\theta}.$$

Lemma 3.

If δ is a permutation-invariant no-data rule, then

$$P(CS|\delta) \geq 1/k$$

Proof.

$$P(CS|\delta) = \sum_{i=1}^k p_i \frac{i}{k} \geq \frac{1}{k} \sum_{i=1}^k p_i = 1/k$$

Q.E.D.

Hence, there are no permutation-invariant no-data rule which can achieve $P^* < 1/k$, showing that Gupta's statement is incorrect for this class.

We also see that the lower bound $1/k$ is achieved by the rule p with $p_1 = 1$.

Now, for any $P^* \geq 1/k$ there exists a no-data rule p with $P(\text{CS}|\underline{p}) = P^*$.

This can be seen as follows. If $1/k \leq P^* < 1/2$, let for example p be given by:

$$p_1 = \frac{k+2-2kP^*}{k}, \quad p_2 = \dots = p_k = \frac{2kP^*-2}{k(k-1)}$$

Then

$$P(\text{CS}|\underline{p}) = P^*.$$

If $P^* \geq 1/2$, let $p_1 + \dots + p_{k-1} = \frac{2(1-P^*)}{k-1}$ and $p_k = 2P^*-1$.

Then

$$P(\text{CS}|\underline{p}) = P^*.$$

Since $E(S|\underline{p}) = kP(\text{CS}|\underline{p})$, it follows from Berger (1979), (under weak regularity conditions) that for each $\frac{1}{k} \leq P^* \leq 1$ there exists a permutation-invariant no-data rule that subject to the condition

$$\inf_{\theta \in \Omega} P_{\theta}(\text{CS}|\delta) \geq P^*$$

is
are minimax for the risk $E_{\theta}(S|\delta)$.

The criterium used by Bechhofer and Santner (1979), which is to choose P^* greater than or equal to $P(\text{CS}|\underline{p}^0)$ where \underline{p}^0 is minimax therefore seems hard to understand.

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